

THE SECOND FENG-RAO NUMBER FOR CODES COMING FROM TELESCOPIC SEMIGROUPS

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ABSTRACT. In this manuscript we show that the second Feng-Rao number of any telescopic numerical semigroup agrees with the multiplicity of the semigroup. To achieve this result we first study the behavior of Apéry sets under gluings of numerical semigroups. These results provide a bound for the second Hamming weight of one-point Algebraic Geometry codes, which improves upon other estimates such as the Griesmer Order Bound.

INTRODUCTION

In coding theory algebraic geometry codes (AG codes for short) are advantageous in that their parameters asymptotically exceed the Gilbert-Varshamov bound (see [HvLP]), and due to Feng and Rao [FR] we know that they may be efficiently decoded. A bound on the number of correctable errors depends on the Feng-Rao distance of the involved Weierstrass semigroup. Even though the original codes are defined over algebraic curves, the construction can avoid the explicit use of algebraic geometry, by means of arrays of codes [KP] and order functions over algebras [HvLP].

We begin with a brief overview of how Feng-Rao distances and numerical semigroups arise in coding theory. Let \mathbb{F}_q be the finite field with q a prime power number of elements. Let R be the affine coordinate ring of a curve over \mathbb{F}_q that is absolutely irreducible, nonsingular and with a single point at infinity. Denote the point at infinity by Q , and let $\mathcal{P} = (P_1, \dots, P_n)$ be a list of (affine) \mathbb{F}_q -rational points on the curve. The evaluation map $\text{ev}_{\mathcal{P}} : R \rightarrow \mathbb{F}_q^n$ is defined as $\text{ev}_{\mathcal{P}}(f) = (f(P_1), \dots, f(P_n))$. Let $v_Q : K \rightarrow \mathbb{Z}$ be the discrete valuation at Q , where K is the quotient ring for R , and define $L(aQ) = \{f \in R \mid v_Q(f) \geq -a\}$. The set $\Gamma := -v_Q(R)$ is a numerical semigroup; see [RG] for the definitions and basic properties of numerical semigroups. For a in Γ , let C_a be the orthogonal linear space of $\text{ev}_{\mathcal{P}}(L(aQ))$ (with respect to the usual dot product). This vector space is called the *one point algebraic code* defined by R , a and \mathcal{P} .

The minimum distance of the code C_a has a lower bound given by the *Feng-Rao distance* [FR] of $a + 1$. The Feng-Rao distance is defined by $\delta_{FR}(a) = \min\{\#D(b) \mid a \leq b, b \in \Gamma\}$, where

2010 *Mathematics Subject Classification.* 20M14, 11T71, 11Y55.

Key words and phrases. Numerical semigroups, telescopic numerical semigroups, free numerical semigroups, gluing of numerical semigroups, Apéry sets, Feng-Rao numbers, Feng-Rao distances, AG codes, generalized Hamming weights.

The first author is supported by the project MTM2015-65764-C3-1-P (MINECO/FEDER).

The second author is supported by the projects MTM2014-55367-P, FQM-343, FQM-5849, and FEDER funds.

The third author is supported by the Fundação para a Ciência e a Tecnologia (Portuguese Foundation for Science and Technology) through the project UID/MAT/00297/2013 (Centro de Matemática e Aplicações).

The fourth author would like to thank Marco D'Anna and the rest of organizers of the INdAM meeting: International meeting on numerical semigroups - Cortona 2014.

The authors would like to thank Vítor Hugo Fernandes for the helpful discussions during the preparation of this paper.

$D(b) = \{c \in \Gamma \mid b - c \in \Gamma\}$ denotes the set of “divisors” of b in Γ (according to the terminology in [AG2]). In order to avoid ambiguity about which semigroup is being used, at times we may write $D_\Gamma(b)$ instead of $D(b)$. Let c and g respectively denote the *conductor* and *genus* of Γ . Then for $a \geq 2c - 1$ we have $\delta_{FR}(a) = a + 1 - 2g$, which is referred to as the Goppa bound. Moreover, one has $\delta_{FR}(a) \geq a + 1 - 2g$ for $a \geq c$.

A natural generalization of the Feng-Rao distance is the following. Take a sequence $a_1 < \dots < a_r$ of r elements in Γ , and define $D(a_1, \dots, a_r) = \bigcup_{i=1}^r D(a_i)$. The r^{th} *generalized Feng-Rao distance* is defined as

$$\delta_{FR}^r(a) = \min\{\#D(a_1, \dots, a_r) \mid a_1, \dots, a_r \in \Gamma, a \leq a_1 < a_2 < \dots < a_r\}.$$

The r^{th} generalized Feng-Rao distance of $a + 1$ turns out to be a lower bound for the r^{th} generalized Hamming weight over the code C_a (see [HP]). We remark that the generalized Hamming weights were introduced independently by Helleseth et al. in [HKM] and Wei in [W], for applications in coding theory and cryptography respectively.

As with the case where $r = 1$, [FM, Theorem 3] shows that for fixed r the asymptotic behavior of δ_{FR}^r is linear. In particular for all $a \geq 2c - 1$ we have $\delta_{FR}^r(a) = a + 1 - 2g + E(\Gamma, r)$ for some constant $E(\Gamma, r)$, known as the r^{th} *Feng-Rao number* of Γ . Furthermore, as in the classical case the inequality $\delta_{FR}^r(a) \geq a + 1 - 2g + E(\Gamma, r)$ holds for $a \geq c$. In [FM, Proposition 5] it is shown that for $g > 0$ and $r \geq 2$, we have $2 \leq E(\Gamma, r) \leq \rho_r$, where ρ_r is the r^{th} smallest element of Γ . Clearly $E(\Gamma, 1) = 0$.

In [DFGL2, FM] it is shown that $E(\Gamma, 2) = \min\{\#\text{Ap}(\Gamma, x) \mid x \in \mathbb{N} \setminus \{0\}\}$, where $\text{Ap}(\Gamma, x) = \{a \in \Gamma \mid a - x \notin \Gamma\}$ is the *Apéry set* of an integer x with respect to Γ . Many of our results are dependent on this alternative formulation. Originally in [A] Apéry only considered the case where x was an element of Γ . Since then natural generalizations have been introduced in [FM] and [GL], which extend the study of Apéry sets allowing for x to be any integer.

Previously the cases where the second Feng-Rao number of a numerical semigroup was known were fairly limited. In [FM] it is shown that when Γ is two generated, the second Feng-Rao number is precisely the *multiplicity* of the semigroup; that is, we have $E(\Gamma, 2) = \rho_2$. Later in [DFGL2] this result was generalized to show that for every r and every two generated numerical semigroup we have $E(\Gamma, r) = \rho_r$. A formula for the r^{th} Feng-Rao number was also calculated for the family of numerical semigroups generated by intervals [DFGL1]. Additionally an expression for $E(\Gamma, 2)$ has been found for the case where Γ is inductive [FG].

It is often the case that properties of two generated numerical semigroups generalize to the class of complete intersection numerical semigroups or more generally to the class of symmetric numerical semigroups. However, even restricting to three generated complete intersection case, there are examples where $E(\Gamma, 2) \neq \rho_2$; see for instance Example 21. A relevant subclass of complete intersection numerical semigroups is the class of telescopic numerical semigroups, which were introduced by Kirfel and Pellikaan in [KP] in order to study Feng-Rao distances. This family contains the set of numerical semigroups associated to irreducible plane curve singularities, which were introduced by Zariski in [Z]. Based on computational evidence it has been suggested for some time that $E(\Gamma, 2) = \rho_2$ whenever Γ is telescopic. In our case we used the GAP [GAP] package `numericalsgps` [DGM] that implements procedures from [AG1] to calculate $E(\Gamma, 2)$ for all telescopic numerical semigroups with genus less than 150. In this paper with the help of auxiliary results on Apéry sets over gluings of numerical semigroups, we show that the second Feng-Rao number of any telescopic numerical semigroup agrees with its multiplicity (Corollary 16). It is our hope that the results from our first section will also prove useful in deriving a formula for $E(\Gamma, 2)$ for more general cases of Γ .

Although Feng-Rao distances and Feng-Rao numbers yield important information about AG codes, it is important to note that the numbers themselves are only dependent on the associated numerical semigroup. Consequently the focus of our research has been to develop tools that allow us to calculate properties of $\text{Ap}(\Gamma, x)$ for a broader class of numerical semigroups. In particular we develop formulas for calculating properties of Apéry sets of gluings on numerical semigroups that may be derived iteratively from the numerical semigroups in the gluing. Some of these formulas generalize a specialized result appearing in [RG, Chapter 8]. The idea of gluing was originally developed to construct curve singularities with particular properties (see for instance [BC]) and was later generalized in [D], and made explicit in [R]. Additionally invariants such as the Frobenius number, conductor, genus type, symmetry and Hilbert series can all be recovered from the original semigroups (see for instance [AGO] and the references therein). In our case by understanding Apéry sets under gluings, we are able to develop tools for calculating the second Feng-Rao number for gluings of numerical semigroups in certain cases. In addition we feel that understanding the construction of Apéry sets under gluings is fundamental to the theory of numerical semigroups, and will thus yield other applications in the future.

The paper is organized as follows: In section 1 we prove fundamental results concerning gluings of numerical semigroups and Apéry sets. Section 2 is devoted to prove the main results of the paper on the second Feng-Rao number for free numerical semigroups, and in particular Corollary 16. Section 3 computes some interesting examples of free semigroups appearing in coding theory, namely the generalized Hermitian semigroups and the Suzuki semigroups. Finally, section 4 applies the previous results to AG codes constructed from the generalized Hermitian curves and the Suzuki curves, obtaining bounds for the second Hamming weight of these codes that are better than the one given by Kirfel-Pellikaan in [KP] and the Griesmer order bound introduced in [DFGL2].

1. GLUINGS AND APÉRY SETS

In order to exploit the formulation of the second Feng-Rao number in terms of Apéry sets, we develop some basic properties of Apéry sets and then show how they behave under gluings. We begin with some standard definitions, which mirror the notation in [RG]. Let \mathbb{N} denote the non-negative integers. A numerical semigroup is a set of the form $\Gamma = \langle n_1, \dots, n_m \rangle := n_1\mathbb{N} + \dots + n_m\mathbb{N}$ where n_1, \dots, n_m are positive integers, such that $\gcd(n_1, \dots, n_m) = 1$. The set $A = \{n_1, \dots, n_m\}$ is said to be the minimal generating set for Γ provided that $\Gamma := \langle A \rangle \neq \langle A \setminus \{n_i\} \rangle$ for any i . Let Γ_1 and Γ_2 be numerical semigroups, and choose $a_1 \in \Gamma_2$ and $a_2 \in \Gamma_1$, such that $\gcd(a_1, a_2) = 1$ and neither a_1 nor a_2 are minimal generators. Then the set $\Gamma = a_1\Gamma_1 + a_2\Gamma_2$ is again a numerical semigroup, referred to as a *gluing* of Γ_1 and Γ_2 .

We can say more than [CGM, Proposition 1] or [FM, Proposition 18].

Lemma 1. *Given a numerical semigroup Γ and $x \in \mathbb{Z}$, we have $\# \text{Ap}(\Gamma, x) = x + \# \text{Ap}(\Gamma, -x)$. Thus $\# \text{Ap}(\Gamma, x) \geq x$ with equality only if $\text{Ap}(\Gamma, -x) = \emptyset$ or equivalently when $x \in \Gamma$.*

Proof. We restrict to the case $x > 0$ and the general case follows by replacing x with $-x$ in the equation $\# \text{Ap}(\Gamma, -x) = \# \text{Ap}(\Gamma, x) - x$.

Let A_i and B_i be the subsets of $\text{Ap}(\Gamma, x)$ and $\text{Ap}(\Gamma, -x)$ respectively whose elements are congruent to i modulo x . Notice that A_i is nonempty, since it contains the smallest element in Γ that is congruent to i modulo x . Let $A_i = \{a_1, a_2, \dots, a_k\}$ with $a_1 < \dots < a_k$. There are no elements in B_i that are greater than or equal to a_k because $a_k + nx - (-x) \in \Gamma$ for all $n \geq 0$. Similarly there are no elements in B_i strictly less than a_1 . Hence if $k = 1$, we have $B_i = \emptyset$. If $k \geq 2$, then for each j with $1 \leq j < k$ there exists $n_j \in \mathbb{N}$ such that $a_j + hx \in \Gamma$ for $0 \leq h \leq n_j$

and $a_j + hx \notin \Gamma$ for $n_j < h < (a_{j+1} - a_j)/x$. It follows that $b_j = a_j + n_j x$ is the unique element of B_i with $a_j \leq b_j < a_{j+1}$. Thus $|A_i| = |B_i| + 1$ and $\# \text{Ap}(\Gamma, -x) = \# \text{Ap}(\Gamma, x) - x$.

If $x \in \Gamma$, then a standard argument shows that $\# \text{Ap}(\Gamma, x) = x$, and hence $\# \text{Ap}(\Gamma, -x)$. Conversely if $x \notin \Gamma$, then $0 \in \text{Ap}(\Gamma, -x) \neq \emptyset$. \square

Lemma 2. *Let Γ be a numerical semigroup. Then we have the following.*

(a) *For $x, y \in \mathbb{Z}$ we have $\# \text{Ap}(\Gamma, x + y) \leq \# \text{Ap}(\Gamma, x) + \# \text{Ap}(\Gamma, y)$.*

(b) *For $g, h \in \Gamma$ we have $\text{Ap}(\Gamma, g + h) = \text{Ap}(\Gamma, g) \cup (g + \text{Ap}(\Gamma, h))$, and the union is disjoint.*

Proof. Let $z \in \text{Ap}(\Gamma, x + y)$. Then $z \in \Gamma$. If $z - x \notin \Gamma$, then $z \in \text{Ap}(\Gamma, x)$. Otherwise $z - x \in \Gamma$ and $z - x - y \notin \Gamma$; hence $z - x \in \text{Ap}(\Gamma, y)$ and $z \in (x + \text{Ap}(\Gamma, y))$. Thus we have the containment $\text{Ap}(\Gamma, x + y) \subseteq \text{Ap}(\Gamma, x) \cup (x + \text{Ap}(\Gamma, y))$. This produces the inequality in (a) and one of the necessary inclusions for (b).

It remains to show the other inclusion of sets for part (b). Let $z' \in \text{Ap}(\Gamma, g)$. Then $z' \in \Gamma$, and $z' - g \notin \Gamma$. Since $h \in \Gamma$ we have $z' - g - h \notin \Gamma$. Thus $z' \in \text{Ap}(\Gamma, g + h)$. Let $z'' \in g + \text{Ap}(\Gamma, h)$. Then $z'' - g \in \Gamma$, and $z'' - g - h \notin \Gamma$. Since $g \in \Gamma$ it follows that $z'' \in \Gamma$, and thus $z'' \in \text{Ap}(\Gamma, g + h)$. \square

Lemma 3. *Let $\Gamma = a_1\Gamma_1 + a_2\Gamma_2$ be a gluing of numerical semigroups. Any integer z can be expressed uniquely as $z = a_1k + a_2\omega$ where $k \in \mathbb{Z}$ is an integer and $\omega \in \text{Ap}(\Gamma_2, a_1)$. Then z is in the semigroup Γ if and only if $k \in \Gamma_1$.*

Proof. Since $\gcd(a_1, a_2) = 1$, and $a_1 \in \Gamma_2$, there exists a unique $\omega \in \text{Ap}(\Gamma_2, a_1)$ such that $z \equiv a_2\omega \pmod{a_1}$. That means that $z = a_1k + a_2\omega$ for a unique $k \in \mathbb{Z}$.

For the second part, it is clear that if $k \in \Gamma_1$, then z is in Γ . Suppose that $z \in \Gamma$, so that $z = a_1x_1 + a_2x_2$ with $x_1 \in \Gamma_1$ and $x_2 \in \Gamma_2$. Then $x_2 \equiv \omega \pmod{a_1}$ and then $x_2 = \omega + na_1$ with $n > 0$. This means that $z = a_1x_1 + a_2x_2 = a_1x_1 + a_2(\omega + na_1) = a_1(x_1 + na_2) + a_2\omega$, so by uniqueness, $k = x_1 + na_2 \in \Gamma_1$. \square

Definition 4. Let Γ be a numerical semigroup and $g \in \Gamma$ an element. Let us write, for $0 \leq i < g$, as $\omega(i)$ the unique element in $\text{Ap}(\Gamma, g)$ such that $\omega(i) \equiv i \pmod{g}$. Then define the cocycle $h_{\Gamma, g} : \mathbb{Z}_g \times \mathbb{Z}_g \rightarrow \mathbb{Z}$ as

$$h_{\Gamma, g}(i, j) = (\omega(i) - \omega(i + j) + \omega(j))/g.$$

Remark 5. From the cocycles $h_{\Gamma, g}(i, j)$ we can recover the elements of $\text{Ap}(\Gamma, g)$ up to congruence. Indeed the formula

$$\omega(i) = \sum_{j=0}^{g-1} h_{\Gamma, g}(j, i)$$

follows easily from the definitions.

Lemma 6. *Let $\Gamma = a_1\Gamma_1 + a_2\Gamma_2$ be a gluing of numerical semigroups. Let $z = a_1k + a_2\omega(i)$ be any integer, with $\omega(i) \in \text{Ap}(\Gamma_2, a_1)$. Then*

$$\text{Ap}(\Gamma, z) = \bigcup_{j=0}^{a_1-1} a_1 \text{Ap}(\Gamma_1, k + a_2 h_{\Gamma_2, a_1}(j - i, i)) + a_2\omega(j),$$

and

$$\text{D}_\Gamma(z) = \bigcup_{j=0}^{a_1-1} a_1 \text{D}_{\Gamma_1}(k - a_2 h_{\Gamma_2, a_1}(i - j, j)) + a_2\omega(j).$$

Proof. Let $s = a_1p + a_2\omega(j)$ an integer. Then

$$\begin{aligned} s - z &= a_1(p - k) + a_2(\omega(j) - \omega(i)) \\ &= a_1(p - k) + a_2(\omega(j - i) - a_1 h_{\Gamma_2, a_1}(j - i, i)) \\ &= a_1(p - k - a_2 h_{\Gamma_2, a_1}(j - i, i)) + a_2\omega(j - i). \end{aligned}$$

Therefore $s \in \text{Ap}(\Gamma, z)$ if and only if $p \in \text{Ap}(\Gamma_1, k + a_2 h_{\Gamma_2, a_1}(j - i, i))$. Similarly

$$\begin{aligned} z - s &= a_1(k - p) + a_2(\omega(i) - \omega(j)) \\ &= a_1(k - p) + a_2(\omega(i - j) - a_1 h_{\Gamma_2, a_1}(i - j, j)) \\ &= a_1(k - a_2 h_{\Gamma_2, a_1}(i - j, j) - p) + a_2 \omega(i - j), \end{aligned}$$

so $s \in D_\Gamma(z)$ if and only if $p \in D_{\Gamma_1}(k - a_2 h_{\Gamma_2, a_1}(i - j, j))$. \square

Proposition 7. *Let $\Gamma = a_1 \Gamma_1 + a_2 \Gamma_2$ be a gluing of numerical semigroups. Let $z = a_1 \alpha + a_2 \omega(i)$ be any integer, with $\omega(i) \in \text{Ap}(\Gamma_2, a_1)$. Let $\beta = \#\{x \in \text{Ap}(\Gamma_2, a_1) \mid x - \omega(i) \notin \Gamma_2\}$. Then*

$$\text{Ap}(\Gamma, z) \geq (a_1 - \beta) \cdot \#\text{Ap}(\Gamma_1, \alpha) + \beta \cdot \#\text{Ap}(\Gamma_1, \alpha + a_2).$$

Proof. We have the following sequence of equalities:

$$\begin{aligned} \omega(j) - \omega(i) \in \Gamma_2 &\iff \omega(j) - \omega(i) \in \text{Ap}(\Gamma_2, a_1) \\ &\iff \omega(j) - \omega(i) = \omega(j - i) \\ &\iff h_{\Gamma_2, a_1}(j - i, i) = (\omega(j - i) - \omega(j) + \omega(i))/a_1 = 0. \end{aligned}$$

Thus $\beta = \#\{j \in \mathbb{Z}_{a_1} \mid h_{\Gamma_2, a_1}(j - i, i) \geq 1\}$. Since $a_2 \in \Gamma_1$ we have $\#\text{Ap}(\Gamma_1, \alpha + a_2 h) \geq \#\text{Ap}(\Gamma_1, \alpha + a_2)$ when $h \geq 1$. This explains the second step below. The first step below comes directly from Lemma 6.

$$\begin{aligned} \#\text{Ap}(\Gamma, z) &= \sum_{j=0}^{a_1-1} \#\text{Ap}(\Gamma_1, \alpha + a_2 h_{\Gamma_2, a_1}(j - i, i)) \\ &\geq (a_1 - \beta) \cdot \#\text{Ap}(\Gamma_1, \alpha) + \beta \cdot \#\text{Ap}(\Gamma_1, \alpha + a_2). \end{aligned} \quad \square$$

Lemma 8. *Let Γ be a numerical semigroup. Choose $x \in \mathbb{Z}$ and $y \in \Gamma$. Then*

$$\text{Ap}(\Gamma, x + y) = \text{Ap}(\Gamma, x) \cup (y + \text{Ap}(\Gamma, x)) \cup (\text{Ap}(\Gamma, y) \cap (x + \text{Ap}(\Gamma, y))).$$

Proof. First we will show that each of the sets in the union on the right is contained on the left side. Suppose that $z \in \text{Ap}(\Gamma, x)$. Since $z - x \notin \Gamma$ we have $z - x - y \notin \Gamma$. Thus $z \in \text{Ap}(\Gamma, x + y)$. Suppose $z' \in (y + \text{Ap}(\Gamma, x))$. Then $z' - y \in \text{Ap}(\Gamma, x)$; hence $z' - y \in \Gamma$ and $z' - y - x \notin \Gamma$. Since $y \in \Gamma$ we have $z' = z' - y + y \in \Gamma$. Thus $z' \in \text{Ap}(\Gamma, x + y)$. Lastly suppose that $z'' \in \text{Ap}(\Gamma, y) \cap (x + \text{Ap}(\Gamma, y))$. Then $z'' \in \text{Ap}(\Gamma, y)$ implies that $z'' \in \Gamma$. Also $z'' - x \in \text{Ap}(\Gamma, y)$ implies that $z'' - x - y \notin \Gamma$. Thus $z'' \in \text{Ap}(\Gamma, x + y)$.

Let $w \in \text{Ap}(\Gamma, x + y)$. Assuming $w \notin \text{Ap}(\Gamma, x) \cup (y + \text{Ap}(\Gamma, x))$ it suffices to show that $w \in (\text{Ap}(\Gamma, y) \cap (x + \text{Ap}(\Gamma, y)))$. Since $w \in \Gamma$, and $w \notin \text{Ap}(\Gamma, x)$ we have that $w - x \in \Gamma$. Thus $w - x \in \text{Ap}(\Gamma, y)$ and $w \in x + \text{Ap}(\Gamma, y)$. Since $w - y \notin \text{Ap}(\Gamma, x)$ and $w - x - y \notin \Gamma$ we have $w - y \notin \Gamma$. Thus $w \in \text{Ap}(\Gamma, y)$ and the result follows. \square

Remark 9. In Lemma 8 we get two special cases:

- (1) If $-x \in \Gamma$, then $\text{Ap}(\Gamma, x + y) = \text{Ap}(\Gamma, y) \cap (x + \text{Ap}(\Gamma, y))$;
- (2) For a set of integers X , write $\Delta(X) = \{x' - x \mid x, x' \in X, x' > x\}$.
If $|x| \notin \Delta(\text{Ap}(\Gamma, y)) \cup \{0\}$, then $\text{Ap}(\Gamma, x + y) = \{0, y\} + \text{Ap}(\Gamma, x)$.

For a gluing $\Gamma = a_1 \Gamma_1 + a_2 \Gamma_2$ the next result produces a simple formula for $\text{Ap}(\Gamma, z)$ in the case where z is divisible by a_1 .

Proposition 10. *Let $\Gamma = a_1 \Gamma_1 + a_2 \Gamma_2$ be a gluing of numerical semigroups. Then given any $z \in \mathbb{Z}$ we have*

$$\text{Ap}(\Gamma, a_1 z) = a_1 \text{Ap}(\Gamma_1, z) + a_2 \text{Ap}(\Gamma_2, a_1).$$

In particular

$$\#\text{Ap}(\Gamma, a_1 z) = a_1 \cdot \#\text{Ap}(\Gamma_1, z).$$

Proof. Using the unique form in Lemma 3 we have $a_1 z = a_1 k + a_2 \omega$ with $\omega = \omega(0) = 0$ and $z = k$. Notice that $\text{h}_{\Gamma_2, a_1}(j - 0, 0) = 0$ for all j ; hence by Lemma 6 we have

$$\text{Ap}(\Gamma, a_1 z) = \bigcup_{j=0}^{a_1-1} a_1 \text{Ap}(\Gamma_1, z) + a_2 \omega(j) = a_1 \text{Ap}(\Gamma_1, z) + a_2 \text{Ap}(\Gamma_2, a_1). \quad \square$$

Notice that by setting $z = a_2$ in Proposition 10 we can recover the previously known special case $\text{Ap}(\Gamma, a_1 a_2) = a_1 \text{Ap}(\Gamma_1, a_2) + a_2 \text{Ap}(\Gamma_2, a_1)$ from [RG, Chapter 8].

Combining Lemma 2 (c) with Proposition 10 we get an expression for $\text{Ap}(\Gamma, z)$ when $z \in \Gamma$.

Corollary 11. *Let $\Gamma = a_1 \Gamma_1 + a_2 \Gamma_2$ be a gluing of numerical semigroups. Let $g_1 \in \Gamma_1$, $g_2 \in \Gamma_2$. and $g = a_1 g_1 + a_2 g_2 \in \Gamma$. Then*

$$\text{Ap}(\Gamma, g) = (a_1 \text{Ap}(\Gamma_1, g_1) + a_2 \text{Ap}(\Gamma_2, a_1)) \cup (a_1 g_1 + a_1 \text{Ap}(\Gamma_1, a_2) + a_2 \text{Ap}(\Gamma_2, g_2)).$$

Since $\#\text{Ap}(\Gamma, g) = g$ for any $g \in \Gamma$, it follows that the union in Corollary 11 is disjoint.

Theorem 12. *Let $\Gamma = a_1 \Gamma_1 + a_2 \mathbb{N}$ be a gluing of numerical semigroups, and let z be an integer. Express z as $z = a_1 \alpha + a_2 \beta$ with $0 \leq \beta < a_1$ and $\alpha \in \mathbb{Z}$. We have the following:*

$$(1) \text{Ap}(\Gamma, z) = (a_1 \text{Ap}(\Gamma_1, \alpha + a_2) + a_2 \{0, \dots, \beta - 1\}) \cup (a_1 \text{Ap}(\Gamma_1, \alpha) + a_2 \{\beta, \dots, a_1 - 1\});$$

$$(2) \text{D}_\Gamma(z) = (a_1 \text{D}_{\Gamma_1}(\alpha) + a_2 \{0, \dots, \beta\}) \cup (a_1 \text{D}_{\Gamma_1}(\alpha + a_2) + a_2 \{\beta + 1, \dots, a_1 - 1\});$$

$$(3) \#\text{Ap}(\Gamma, z) = \beta \cdot \#\text{Ap}(\Gamma_1, \alpha + a_2) + (a_1 - \beta) \cdot \#\text{Ap}(\Gamma_1, \alpha);$$

$$(4) \#\text{D}_\Gamma(z) = (\beta + 1) \#\text{D}_{\Gamma_1}(\alpha) + (a_1 - \beta - 1) \#\text{D}_{\Gamma_1}(\alpha - a_2).$$

Proof. Notice that $\text{Ap}(\mathbb{N}, a_1) = \{0, 1, \dots, a_1 - 1\}$ and $\omega(j) = j$ for $0 \leq j < a_1$. When $j < \beta$ we have $\text{h}_{\mathbb{N}, a_1}(j - \beta, \beta) = 1$, and when $j \geq \beta$ we have $\text{h}_{\mathbb{N}, a_1}(j - \beta, \beta) = 0$. By Lemma 6 we have

$$\begin{aligned} \text{Ap}(\Gamma, z) &= \bigcup_{j=0}^{a_1-1} a_1 \text{Ap}(\Gamma_1, \alpha + a_2 \text{h}_{\Gamma_2, a_1}(j - \beta, \beta)) + a_2 j \\ &= (a_1 \text{Ap}(\Gamma_1, \alpha + a_2) + a_2 \{0, \dots, \beta - 1\}) \cup (a_1 \text{Ap}(\Gamma_1, \alpha) + a_2 \{\beta, \dots, a_1 - 1\}). \end{aligned}$$

When $j \leq \beta$ we have $\text{h}_{\mathbb{N}, a_1}(\beta - j, j) = 0$, and when $j > \beta$ we have $\text{h}_{\mathbb{N}, a_1}(\beta - j, j) = 1$. By Lemma 6 we have

$$\begin{aligned} \text{D}_\Gamma(z) &= \bigcup_{j=0}^{a_1-1} a_1 \text{D}_{\Gamma_1}(\alpha - a_2 \text{h}_{\Gamma_2, a_1}(\beta - j, j)) + a_2 j \\ &= (a_1 \text{D}_{\Gamma_1}(\alpha) + a_2 \{0, \dots, \beta\}) \cup (a_1 \text{D}_{\Gamma_1}(\alpha - a_2) + a_2 \{\beta + 1, \dots, a_1 - 1\}). \end{aligned}$$

$$\text{Hence } \#\text{D}_\Gamma(z) = (\beta + 1) \#\text{D}_{\Gamma_1}(\alpha) + (a_1 - \beta - 1) \#\text{D}_{\Gamma_1}(\alpha - a_2). \quad \square$$

Notices that (4) in Theorem 12 can be sharpened a bit more. If $\alpha \in \text{Ap}(\Gamma_1, a_2)$, then $\text{D}_{\Gamma_1}(\alpha - a_2)$ is empty. If to the contrary, $\alpha - a_2 \in \Gamma$, then $\text{D}_{\Gamma_1}(\alpha) = \text{D}_{\Gamma_1}(\alpha - a_2) \cup \text{D}_{\Gamma_1}(\alpha) \setminus \text{D}_{\Gamma_1}(\alpha - a_2)$. And the cardinality of $\text{D}_{\Gamma_1}(\alpha - a_2 + a_2) \setminus \text{D}_{\Gamma_1}(\alpha - a_2)$ is precisely that of $\text{Ap}(\Gamma_1, a_2)$ ([DFGL2, Proposition 11]). Hence $\#\text{D}_\Gamma(z) = (\beta + 1) \#\text{D}_{\Gamma_1}(\alpha) + (a_1 - \beta - 1) \#\text{D}_{\Gamma_1}(\alpha - a_2) = (\beta + 1)(\#\text{D}_{\Gamma_1}(\alpha - a_2) + a_2) + (a_1 - \beta - 1) \#\text{D}_{\Gamma_1}(\alpha - a_2) = a_1 \#\text{D}_{\Gamma_1}(\alpha - a_2) + (\beta + 1)a_2$. By putting all this together, we obtain.

$$(1) \quad \#\text{D}(z) = \begin{cases} (\beta + 1) \#\text{D}_{\Gamma_1}(\alpha), & \text{if } \alpha \in \text{Ap}(\Gamma, a_2), \\ a_1 \#\text{D}_{\Gamma_1}(\alpha - a_2) + (\beta + 1)a_2, & \text{otherwise.} \end{cases}$$

Remark 13. Suppose $\Gamma = \langle a, b \rangle$ and $z = ua + vb$ with $0 \leq u < b$. Then we may apply Theorem 12 with $\Gamma_1 = \mathbb{N}$, $a_1 = b$, $a_2 = a$, $\beta = u$ and $\alpha = v$. Thus

$$\# \text{Ap}(\Gamma, z) = (b - u) \max\{v, 0\} + u \max\{v + a, 0\}.$$

This yields the same result as [DFGL2, Theorem 14].

2. THE SECOND FENG-RAO NUMBER FOR FREE NUMERICAL SEMIGROUPS

Recall that the second Feng-Rao number of a numerical semigroup Γ can be computed as

$$\begin{aligned} E(\Gamma, 2) &= \min\{\# \text{Ap}(\Gamma, z) \mid z \in \mathbb{N} \setminus \{0\}\} \\ &= \min\{\# \text{Ap}(\Gamma, z) \mid z \in \{1, \dots, m(\Gamma)\}\}. \end{aligned}$$

To realize the second equality above notice that for $x \geq m(\Gamma)$ we have $\# \text{Ap}(\Gamma, x) \geq x \geq m(\Gamma)$.

Remark 14. Let Γ be a numerical semigroup. Then $E(\Gamma, 2) = m(\Gamma)$ if and only if there exists $y \geq m(\Gamma)$ such that $\# \text{Ap}(\Gamma, z) \geq m(\Gamma) \lceil \frac{z}{y} \rceil$ for all $z \in \mathbb{Z}$. If this is the case, then y can be chosen so that it is at most $2m(\Gamma) - 1$.

Clearly if $\# \text{Ap}(\Gamma, z) \geq m(\Gamma) \lceil \frac{z}{y} \rceil$ with $y > 0$, then $\# \text{Ap}(\Gamma, z) \geq m(\Gamma)$ for $z > 0$; hence $E(\Gamma, 2) = m(\Gamma)$. Let $z = nm(\Gamma) + r$ with $0 \leq r < m(\Gamma)$. Notice that for $n \geq 1$ we always have $n \geq \lceil \frac{nm(\Gamma)+r}{2m(\Gamma)-1} \rceil$; hence $\# \text{Ap}(\Gamma, z) \geq nm(\Gamma) + r \geq m(\Gamma) \lceil \frac{nm(\Gamma)+r}{2m(\Gamma)-1} \rceil$. Therefore supposing $E(\Gamma, 2) = m(\Gamma)$ we have $\# \text{Ap}(\Gamma, z) \geq m(\Gamma) \lceil \frac{z}{2m(\Gamma)-1} \rceil$ for all $z \in \mathbb{Z}$. Also notice that if the statement is true for a certain y value then it is true for all $y' > y$. If $0 < y < m(\Gamma)$, then $\# \text{Ap}(\Gamma, m(\Gamma)) = m(\Gamma) < m(\Gamma) \lceil \frac{m(\Gamma)}{y} \rceil$ would not work; hence we must have $y \geq m(\Gamma)$.

Theorem 15. Let $\Gamma = a_1\Gamma_1 + a_2\Gamma_2$ be a gluing of numerical semigroups. Suppose there exists $y \geq m(\Gamma)$ such that $\# \text{Ap}(\Gamma_1, z) \geq m(\Gamma_1) \lceil \frac{z}{y} \rceil$ for all $z \in \mathbb{Z}$. If $a_2 > a_1y$, then $\# \text{Ap}(\Gamma, z) \geq m(\Gamma) \lceil \frac{z}{a_2} \rceil$ for all $z \in \mathbb{Z}$.

Proof. If z is less than or equal to zero the result is clear. Let z be a positive integer and write $z = a_1k + a_2\omega(i)$, where $\omega(j)$ denotes the element of $\text{Ap}(\Gamma_2, a_1)$ which is congruent to j modulo a_1 . For ease of notation set $h_j = h_{\Gamma_2, a_1}(j - i, i)$. As in Lemma 5 note that $\sum_{j=0}^{a_1-1} h_j = \omega(i)$. First suppose $k \geq 0$. Lemma 6 gives us the first step below.

$$\begin{aligned} \# \text{Ap}(\Gamma, z) &= \sum_{j=0}^{a_1-1} \# \text{Ap}(\Gamma_1, k + a_2h_j) \\ &\geq \sum_{j=0}^{a_1-1} m(\Gamma_1) \lceil \frac{k+a_2h_j}{y} \rceil \\ &\geq \sum_{j=0}^{a_1-1} m(\Gamma_1) \lceil \frac{k+a_1yh_j}{y} \rceil \\ &= \sum_{j=0}^{a_1-1} m(\Gamma_1) (\lceil \frac{k}{y} \rceil + a_1h_j) \\ &= a_1m(\Gamma_1) \lceil \frac{k}{y} \rceil + a_1m(\Gamma_1)\omega(i) \\ &= a_1m(\Gamma_1) \lceil \frac{k}{y} + \omega(i) \rceil \\ &= m(\Gamma) \lceil \frac{a_1k}{a_1y} + \frac{a_2\omega(i)}{a_2} \rceil \\ &\geq m(\Gamma) \lceil \frac{a_1k}{a_2} + \frac{a_2\omega(i)}{a_2} \rceil \\ &\geq m(\Gamma) \lceil \frac{z}{a_2} \rceil. \end{aligned}$$

Next suppose $k < 0$. Set $c = a_2 - a_1y$. From the equation $z = a_1k + a_2\omega(i)$ we get $\lceil \frac{z}{a_2} \rceil = \omega(i) - \lfloor \frac{-a_1k}{a_2} \rfloor$, which gives us the last step below. Since $0 < c < a_2$ when $h_j \geq \frac{-k}{a_2}$ we

have $\frac{-k}{a_2} \geq \frac{-k-h_jc}{a_2-c}$, which explains the eighth step below.

$$\begin{aligned}
\#\text{Ap}(\Gamma, z) &= \sum_{j=0}^{a_1-1} \#\text{Ap}(\Gamma_1, k + a_2 h_j) \\
&\geq \sum_{h_j \geq \frac{-k}{a_2}} \#\text{Ap}(\Gamma_1, k + a_2 h_j) \\
&\geq m(\Gamma_1) \sum_{h_j \geq \frac{-k}{a_2}} \left\lceil \frac{k+a_2 h_j}{y} \right\rceil \\
&= m(\Gamma_1) \sum_{h_j \geq \frac{-k}{a_2}} \left\lceil \frac{k+h_j(c+a_1 y)}{y} \right\rceil \\
&= m(\Gamma_1) \sum_{h_j \geq \frac{-k}{a_2}} \left(a_1 h_j + \left\lceil \frac{k+h_j c}{y} \right\rceil \right) \\
&= m(\Gamma_1) \sum_{h_j \geq \frac{-k}{a_2}} \left(a_1 h_j - \left\lfloor \frac{-k-h_j c}{y} \right\rfloor \right) \\
&= m(\Gamma_1) \sum_{h_j \geq \frac{-k}{a_2}} \left(a_1 h_j - \left\lfloor a_1 \frac{-k-h_j c}{a_2-c} \right\rfloor \right) \\
&\geq m(\Gamma_1) \sum_{h_j \geq \frac{-k}{a_2}} \left(a_1 h_j - \left\lfloor a_1 \frac{-k}{a_2} \right\rfloor \right) \\
&= m(\Gamma_1) \left(\sum_{j=0}^{a_1-1} a_1 h_j - \sum_{h_j < \frac{-k}{a_2}} a_1 h_j - \sum_{h_j \geq \frac{-k}{a_2}} \left\lfloor \frac{-a_1 k}{a_2} \right\rfloor \right) \\
&\geq m(\Gamma_1) \left(a_1 \omega(i) - \sum_{h_j < \frac{-k}{a_2}} \left\lfloor \frac{-a_1 k}{a_2} \right\rfloor - \sum_{h_j \geq \frac{-k}{a_2}} \left\lfloor \frac{-a_1 k}{a_2} \right\rfloor \right) \\
&= m(\Gamma_1) \left(a_1 \omega(i) - a_1 \left\lfloor \frac{-a_1 k}{a_2} \right\rfloor \right) \\
&= m(\Gamma) \left\lceil \frac{z}{a_2} \right\rceil. \quad \square
\end{aligned}$$

We say that Γ is *free* if either Γ is \mathbb{N} or it is the a gluing of a free numerical semigroup with \mathbb{N} (these semigroups were introduced by [BC]). Analogously, a numerical semigroup Γ is *telescopic* if either Γ is \mathbb{N} or $\Gamma = a_1 \Gamma_1 + a_2 \mathbb{N}$ is a gluing with Γ_1 telescopic and $a_2 > a_1 n$ for each minimal generator n of Γ_1 .

Corollary 16. *Let $\Gamma = \langle n_1, n_2, \dots, n_m \rangle$ be telescopic with $n_1 < \dots < n_m$. Given $z \in \mathbb{N}$ we have $\#\text{Ap}(\Gamma, z) \geq n_1 \lceil \frac{z}{n_m} \rceil$; hence $E(\Gamma, 2) = n_1$.*

Proof. We proceed by induction on m . For $\Gamma = \mathbb{N}$, the result is obviously true. Now suppose $\Gamma = \langle n_1, n_2, \dots, n_m \rangle$ is telescopic, that is $\Gamma = d\Gamma_1 + n_m \mathbb{N}$ where $d = \gcd(n_1, \dots, n_{m-1})$, $\Gamma_1 = \langle n_1/d, \dots, n_{m-1}/d \rangle$ is telescopic, and $n_m > n_{m-1}$. By induction hypothesis $\#\text{Ap}(\Gamma_1, z) \geq m(\Gamma_1) \lceil \frac{z}{n_{m-1}/d} \rceil$ for all $z \in \mathbb{Z}$, and so by Theorem 15 we get $\#\text{Ap}(\Gamma, z) \geq m(\Gamma) \lceil \frac{z}{n_m} \rceil$. \square

If we take $y = 1$ and $\Gamma_1 = \mathbb{N}$ in Theorem 15, we obtain the following consequence.

Corollary 17. *Let $\Gamma = a_1 \mathbb{N} + a_2 \Gamma_2$ be a gluing of numerical semigroups. If $a_2 > a_1$, then $\#\text{Ap}(\Gamma, z) \geq a_1 \lceil \frac{z}{a_2} \rceil$ for all $z \in \mathbb{Z}$; hence $E(\Gamma, 2) = a_1 = m(\Gamma)$.*

As we mentioned above, if Γ has embedding dimension two, then $E(\Gamma, r)$ corresponds with the r^{th} smallest entry of Γ . However this is no longer true for $r > 2$ in a telescopic numerical semigroup.

Example 18. Let $\Gamma = \langle 6, 10, 11 \rangle$. Then Γ is a telescopic numerical semigroup,

$$\Gamma = \{0, 6, 10, 11, 12, 16, 17, 18, 20, 21, 22, 23, 24, 26, \rightarrow\},$$

and $E(\Gamma, 3) = 9 < 10$.

By [RG, Chapter 8] a numerical semigroup $\Gamma \neq \mathbb{N}$ is free if and only if there exists an arrangement of its minimal generators (n_1, \dots, n_e) such that for $d_i := \gcd(n_1, \dots, n_{i-1})$ and $c_i := d_i/d_{i+1}$ we have $d_i > d_{i+1}$ and $c_i n_i \in \langle n_1, \dots, n_{i-1} \rangle$ for $i \in \{2, \dots, e\}$. In this case we will say that Γ is free with respect to the arrangement (n_1, \dots, n_e) .

Lemma 19. *Let $\Gamma = a_1\Gamma_1 + a_2\mathbb{N}$ be a gluing of numerical semigroups and z a positive integer. We may write $z = a_1\alpha + a_2\beta$ with $0 \leq \beta < a_1$ and $\alpha \in \mathbb{Z}$. Then we have the following:*

$$\#\text{Ap}(\Gamma, z) > (a_1 - \beta) \left(\#\text{Ap}(\Gamma_1, \alpha) + \frac{a_2\beta}{a_1} \right).$$

- If $\alpha > 0$ (this includes the case $\beta = 0$), then $\#\text{Ap}(\Gamma, z) \geq a_1 E(\Gamma_1, 2)$.
- If $\beta \in \{2, 3, \dots, a_1 - 2\}$, then $\#\text{Ap}(\Gamma, z) > a_2$.
- If $\beta \in \{1, a_1 - 1\}$, then $\#\text{Ap}(\Gamma, z) > \frac{a_1 - 1}{a_1} a_2$.

$$\text{Thus } E(\Gamma, 2) \geq \min \left\{ a_1 E(\Gamma_1, 2), \frac{(a_1 - 1)a_2}{a_1} \right\}.$$

Proof. By Theorem 12, we have the first step below.

$$\begin{aligned} \#\text{Ap}(\Gamma, z) &= (a_1 - \beta) \#\text{Ap}(\Gamma_1, \alpha) + \beta \cdot \#\text{Ap}(\Gamma_1, \alpha + a_2) \\ &\geq (a_1 - \beta) \#\text{Ap}(\Gamma_1, \alpha) + \beta(\alpha + a_2) \\ &> (a_1 - \beta) \#\text{Ap}(\Gamma_1, \alpha) + \beta \left(-\frac{a_2\beta}{a_1} + a_2 \right) \\ &= (a_1 - \beta) \left(\#\text{Ap}(\Gamma_1, \alpha) + \frac{a_2\beta}{a_1} \right). \end{aligned}$$

The second step above follows from Lemma 1. The third step follows from the assumption $z = a_1\alpha + a_2\beta > 0$, and the last step is elementary algebra.

If $\alpha > 0$, then $\#\text{Ap}(\Gamma_1, \alpha + a_2) \geq \#\text{Ap}(\Gamma_1, \alpha) \geq E(\Gamma_1, 2)$; hence

$$\#\text{Ap}(\Gamma, z) \geq (a_1 - \beta)E(\Gamma_1, 2) + \beta E(\Gamma_1, 2) = a_1 E(\Gamma_1, 2).$$

If $\beta \in \{2, 3, \dots, a_1 - 2\}$, then $a_1 \geq 4$. Since $\beta(a_1 - \beta)$ has a unique local maximum at $\beta = \frac{a_1}{2}$ we have $\beta(a_1 - \beta) \geq 2(a_1 - 2)$. Since $a_1 \geq 4$, we have $2(a_1 - 2) \geq a_1$. Thus

$$\#\text{Ap}(\Gamma, z) > \beta(a_1 - \beta) \frac{a_2}{a_1} \geq a_1 \frac{a_2}{a_1} = a_2.$$

If $\beta \in \{1, a_1 - 1\}$, then $(a_1 - \beta)\beta = a_1 - 1$. Thus $\#\text{Ap}(\Gamma, z) > (a_1 - \beta) \frac{a_2\beta}{a_1} = \frac{a_1 - 1}{a_1} a_2$. \square

Theorem 20. *Let Γ be a free numerical semigroup with respect to the arrangement (n_1, \dots, n_e) of its minimal generators. If $\frac{c_i - 1}{c_i} n_i \geq n_1$ for $i \in \{3, \dots, e\}$, then $E(\Gamma, 2) = m(\Gamma)$.*

Proof. For $e = 1$ we have $\Gamma = \mathbb{N}$. For $e = 2$ the result is known to be true; see for instance [DFGL2]. Suppose that $e \geq 3$ and that the result is true for the case with $e - 1$ generators. Then $E(\Gamma', 2) = m(\Gamma') = \frac{n_1}{d_e}$ since the sequence of c_i 's for Γ' is that of Γ after removing the last one. Note that $c_e = d_e$. We have that $\Gamma = c_e\Gamma' + n_e\mathbb{N}$. We have $E(\Gamma, 2) \geq \min\{c_e E(\Gamma', 2), \frac{(c_e - 1)n_e}{c_e}\}$ by Lemma 19. Note that $c_e E(\Gamma', 2) = n_1 = m(\Gamma)$ and by hypothesis $\frac{(c_e - 1)n_e}{c_e} \geq n_1 = m(\Gamma)$. Thus $E(\Gamma, 2) \geq m(\Gamma)$. Since we always have $E(\Gamma, 2) \leq m(\Gamma)$, the result follows. \square

If we remove the condition $\frac{c_i - 1}{c_i} n_i \geq n_1$, it may happen that $E(\Gamma, 2) < m(\Gamma)$.

Example 21. Let $\Gamma = \langle 4, 5, 6 \rangle$. Then $\Gamma = 2\langle 2, 3 \rangle + 5\mathbb{N}$ is a gluing of numerical semigroups. Notice that Γ is free but not telescopic. Also $\frac{1}{2}5 = \frac{c_3 - 1}{c_3} n_3 < n_1 = 4$. In this case we have

$$E(\Gamma, 2) = \#\text{Ap}(\Gamma, 1) = |\{0, 4, 8\}| = 3 < 4 = m(\Gamma).$$

Remark 22. It is possible to apply Theorem 12 iteratively to find explicit formulas for $\#\text{Ap}(\Gamma, z)$ when Γ is a free numerical semigroups with more than two generators. However, the complexity of these formulas tends to increase exponentially with the number of generators. Let $M(x) = \max\{x, 0\}$

For example suppose $\Gamma = \langle a, b, c \rangle$ is a three generated complete intersection numerical semigroup (and thus free). Then up to a permutation of the generators we may write $a = \sigma x$, $b = \sigma y$ and $c = c_x x + c_y y$ where $\sigma, x, y > 1$, $x > c_y \geq 0$, $c_x \geq 0$, $\gcd(x, y) = 1$ and $\gcd(\sigma, c) = 1$. Given any integer z we may write z uniquely as $z = z_\sigma \sigma + z_c c$ with $0 \leq z_c < \sigma$, $z_\sigma = z_x x + z_y y$ and $0 \leq z_y < x$.

If $z_y + c_y < x$, then $\#\text{Ap}(\Gamma, z)$ may be expressed as

$$(\sigma - z_c)((x - z_y)M(z_x) + z_y M(z_x + y)) \\ + z_c((x - z_y - c_y)M(z_x + c_x) + (z_y + c_y)M(z_x + c_x + y)).$$

If $z_y + c_y \geq x$, then $\#\text{Ap}(\Gamma, z)$ may be expressed as

$$(\sigma - z_c)((x - z_y)M(z_x) + z_y M(z_x + y)) \\ + z_c((2x - z_y - c_y)M(2z_x + y) + (z_y + c_y - x)M(z_x + c_x + 2y)).$$

3. SOME INTERESTING EXAMPLES

In this section, we apply the above calculations to some interesting examples, coming from the theory of AG codes.

3.1. Generalized Hermitian semigroups. As a generalization of classical Hermitian semigroups, $\langle q, q+1 \rangle$ with q an integer greater than 2, we consider generalized Hermitian semigroups. These are three generated semigroups depending on two parameters: q as above and an integer $r > 2$. The generalized Hermitian curve χ_r with parameters q and r , is defined over \mathbb{F}_{q^r} , by the equation

$$Y^{q^{r-1}} + \dots + Y^q + Y = X^{1+q} + \dots + X^{q^{r-2}+q^{r-1}},$$

and has $q^{2r-1} + 1$ rational points over \mathbb{F}_{q^r} . Its Weierstrass semigroup at the unique pole of X is precisely $H_{q,r} = \langle q^{r-1}, q^{r-1} + q^{r-2}, q^r + 1 \rangle$. For background on this topic see [MST]. Note that χ_2 is the classical Hermitian curve.

Observe that $H_{q,r}$ is the gluing of $\langle q, q+1 \rangle$ and \mathbb{N} given by $H_{q,r} = q^{r-2}\langle q, q+1 \rangle + (q^r + 1)\mathbb{N}$. Thus these are telescopic numerical semigroups, and Corollary 16 yields $E(H_{q,r}, 2) = q^{r-1}$.

3.2. Generalized Suzuki numerical semigroups. Given positive integers p and n , define the generalized Suzuki numerical semigroup as

$$S_{p,n} = \langle p^{2n+1}, p^{2n+1} + p^n, p^{2n+1} + p^{n+1}, p^{2n+1} + p^{n+1} + 1 \rangle$$

(see [M, MR]). The cases where $p = 2$ and n varies are called Suzuki numerical semigroups. Suzuki numerical semigroups come from the Suzuki curve χ^n defined by the equation

$$Y^{2^{2n+1}} + Y = X^{2^{2n}}(X^{2^{2n+1}} + X).$$

These curves have numerous rational points over $\mathbb{F}_{2^{2n+1}}$, which makes them useful for coding theory purposes. By defining $\Gamma_1(p, n) := \langle p^{n+1}, p^{n+1} + 1, p^{n+1} + p \rangle$, we obtain the gluing

$$S_{p,n} = p^n \Gamma_1(p, n) + (p^{2n+1} + p^{n+1} + 1)\mathbb{N}.$$

Notice that $\Gamma_1(p, n) = p\langle p^n, p^n + 1 \rangle + (p^{n+1} + 1)\mathbb{N}$ is also a gluing but is not telescopic.

Lemma 23. $E(\Gamma_1(p, n), 2) = p^{n+1} - p^n + 1$.

Proof. Notice that

$$pE(\langle p^n, p^n + 1 \rangle, 2) = p^{n+1} > p^{n+1} - p^n + 1 - \frac{1}{p} = \frac{(p-1)(p^{n+1} + 1)}{p}.$$

Thus by Lemma 19 we have $E(\Gamma_1(p, n), 2) \geq p^{n+1} - p^n + 1 - \frac{1}{p}$.

By rounding up to the nearest integer we obtain $E(\Gamma_1(p, n), 2) \geq p^{n+1} - p^n + 1$.

Now write $1 = p(-p^n) + (p^{n+1} + 1)$. Using the notation in Theorem 12 for $z = 1$ we have $\alpha = -p^n$, $\beta = 1$, $a_1 = p$ and $a_2 = p^{n+1} + 1$. Thus

$$\begin{aligned} \# \text{Ap}(\Gamma_1(p, n), 1) &= \# \text{Ap}(\langle p^n, p^n + 1 \rangle, -p^n + p^{n+1} + 1) + (p-1) \# \text{Ap}(\langle p^n, p^n + 1 \rangle, -p^n) \\ &= p^{n+1} - p^n + 1, \end{aligned}$$

and the result follows. \square

Theorem 24. $E(S_{p,n}, 2) = p^{2n+1} - p^{2n} + p^n$.

Proof. Again we can compute

$$p^n E(\Gamma_1(p, n), 2) = p^{2n+1} - p^{2n} + p^n.$$

and

$$\frac{(p^n - 1)(p^{2n+1} + p^{n+1} + 1)}{p^n} = p^{2n+1} - p + 1 - \frac{1}{p^n}.$$

Since the first of these values is the smallest, Lemma 19 gives us the inequality

$$E(S_{p,n}, 2) \geq p^{2n+1} - p^{2n} + p^n.$$

From Proposition 10 we obtain $\# \text{Ap}(S_{p,n}, p^n) = p^n \# \text{Ap}(\Gamma_1(p, n), 1) = p^{2n+1} - p^{2n} + p^n$, and the result follows. \square

4. APPLICATION TO AG CODES

Corollary 16 for telescopic semigroups as well as Theorem 24 for Suzuki semigroups provide us with estimates for the second Hamming weight of codes in the array of AG codes corresponding to these numerical semigroups (see [KP]). We recall briefly the definition of the generalized (Hamming) weights. In fact the support of a linear code C is defined as

$$\text{supp}(C) := \{i \mid c_i \neq 0 \text{ for some } \mathbf{c} \in C\},$$

and the r^{th} generalized weight of C is then

$$d_r(C) := \min\{\#\text{supp}(C') \mid C' \preceq C \text{ with } \dim(C') = r\},$$

where $C' \preceq C$ denotes that C' is a linear subcode of C . Observe that the above definition only makes sense if $r \leq k$, where k is the dimension of C .

Let C_a be a code in an array of codes as in [KP] with associated semigroup Γ . For example, C_a may be a one-point AG code associated to a divisor of the form $G = aP$. In this case Γ would be the Weierstrass semigroup of the underlying curve at P , as explained in the introduction. We will consider cases where Γ is telescopic as in Corollary 16 or free with assumptions as in Theorem 20, so that the second Feng-Rao number equals n_1 , the multiplicity of the semigroup. We will also consider the case of the Suzuki semigroups.

Since free semigroups are symmetric (see [RG, Chapter 8]), the bound for the second Hamming weight given by the second Feng-Rao number gives the exact value of the second Feng-Rao distance for half of the elements in the interval $[c, 2c - 1]$, where $c = 2g$ and $a = 2g - 1 + \rho$ with $\rho \in \Gamma \setminus \{0\}$ (see [FM]).

a	12	13	14	15	16	17	18	19	20	21	22	23
δ_{a+2}	4	4	6	6	8	8	9	10	12	12	13	13
$GOB(a+1) := \delta_{a+1} + \lceil \delta_{a+1}/q^r \rceil$	5	5	5	7	7	9	9	11	12	14	14	15
$a+1-2g+E_2$	6	7	8	9	10	11	12	13	14	15	16	17
δ_{a+1}^2	6	8	8	9	10	11	12	13	14	15	16	17

TABLE 1. Results for the case $H_{2,3}$.

Let $a \geq c$. Corollary 16 implies that

$$(2) \quad d_2(C_a) \geq \delta_{FR}^2(a+1) \geq a+2-2g+E_2 = a+2-2g+n_1.$$

In contrast following [KP, Theorem 2.8], one deduces that

$$(3) \quad d_2(C_a) \geq \delta_{FR}(a+2) \geq a+3-2g.$$

Since $n_1 > 1$ the bound obtained from Corollary 16 is an improvement.

Finally, the Griesmer order bound (introduced in [DFGL2]) for the case $r = 2$ shows that

$$(4) \quad d_2(C_a) \geq GOB(a+1) := \delta_{FR}(a+1) + \left\lceil \frac{\delta_{FR}(a+1)}{q} \right\rceil,$$

where the code is defined over the finite field \mathbb{F}_q . In the case of codes constructed from generalized Hermitian semigroups, the underlying field for the corresponding codes would be \mathbb{F}_{q^r} . Consequently we would divide by q^r instead q in the formula above.

In the following examples we compare our bound (2) for the second Hamming weight with (3) and (4). We will denote $\delta_{FR}(a) = \delta_a$ and $\delta_{FR}^2(a) = \delta_a^2$ for simplicity.

Example 25. Consider the generalized Hermitian semigroups as in section 3.1, that is

$$\Gamma = H_{q,r} = \langle q^{r-1}, q^{r-1} + q^{r-2}, q^r + 1 \rangle$$

where q is the power of a prime number and $r \geq 3$ (note that if $r = 2$ the semigroup is Hermitian and is generated by only two elements).

First, consider the case $q = 2$ and $r = 3$. Here $\Gamma = \langle 4, 6, 9 \rangle$ with genus $g = 6$. Table 1 shows the index a of the code C_a together with four bounds for the second Hamming weight, namely bounds (3), (4), (2) and finally the actual second Feng-Rao distance δ_{a+1}^2 . We observe that our bound (2) is always better than Kirfel-Pellikaan bound (3) and the Griesmer order bound (4). Additionally in this case, all of the values but one for our bound (2) coincide exactly with the value of the second Feng-Rao distance.

Next, consider the case $\Gamma = H_{2,4} = \langle 8, 12, 17 \rangle$ with genus $g = 28$. Again for the majority of the values of a our bound (2) coincides exactly with the value of the second Feng-Rao distance. Additionally for all values of a our bound improves upon the bounds (3) and (4). This is apparent for the values displayed in Table 2. Note that for the sake of brevity we omit displaying most of the columns in the table.

We obtained similar results for higher values of q and r with the aid of GAP.

In the next example we use Theorem 24 to obtain the bound

$$(5) \quad d_2(C_a) \geq \delta_{FR}^2(a+1) \geq a+2-2g+E_2 = a+2-2g+(p^{2n+1}-p^{2n}+p^n).$$

These codes are defined solely in the case where characteristic $p = 2$; see [M, MR] for details.

a	56	57	58	59	60	61	62	...	104	105	106	107	108	109	110	111
δ_{a+2}	8	8	8	8	8	8	12	...	51	52	53	54	56	56	57	57
$GOB(a+1)$	9	9	9	9	9	9	9	...	54	55	56	57	58	60	60	61
$a+1-2g+E_2$	10	11	12	13	14	15	16	...	58	59	60	61	62	63	64	65
δ_{a+1}^2	12	12	12	16	16	16	16	...	58	59	60	61	62	63	64	65

TABLE 2. Results for the case $H_{2,4}$.

a	28	29	30	31	32	33	34	35	36	37	38	39	40	41
δ_{a+2}	8	8	8	8	8	8	10	10	12	12	13	16	16	16
$GOB(a+1)$	7	9	9	9	9	9	9	12	12	14	14	15	18	18
$a+1-2g+E_2$	8	9	10	11	12	13	14	15	16	17	18	19	20	21
δ_{a+1}^2	10	11	12	12	12	14	14	16	16	17	18	19	20	22
a	42	43	44	45	46	47	48	49	50	51	52	53	54	55
δ_{a+2}	18	18	20	20	21	22	23	24	25	26	28	28	29	29
$GOB(a+1)$	18	21	21	23	23	24	25	26	27	29	30	32	32	33
$a+1-2g+E_2$	22	23	24	25	26	27	28	29	30	31	32	33	34	35
δ_{a+1}^2	22	24	24	25	26	27	28	29	30	31	32	33	34	35

TABLE 3. Results for the Suzuki semigroup $S_{2,1}$.

Example 26. Consider now the Suzuki semigroup from Section 3.2 for the case $p = 2$

$$\Gamma = S_{2,n} = \langle 2^{2n+1}, 2^{2n+1} + 2^n, 2^{2n+1} + 2^{n+1}, 2^{2n+1} + 2^{n+1} + 1 \rangle.$$

As noted earlier this semigroup is not telescopic. According to Theorem 24, the second Feng-Rao number is $E(\Gamma, 2) = 2^{2n+1} - 2^{2n} + 2^n$.

Consider the classical Suzuki curve with $n = 1$, with Weierstrass semigroup $\Gamma = \langle 8, 10, 12, 13 \rangle$. This semigroup is free with genus $g = 14$, conductor $c = 28$, and second Feng-Rao number $E_2 = 6$. As is apparent in Table 3, our bound is better than both (3) and (4). Additionally we see that our bound equals the second Feng-Rao distance for a majority of the values. Note that in computing the Griesmer order bound we divide by the size of the finite field i.e. 2^{2n+1} .

We also performed computations for higher n , obtaining similar results with large tables.

In conclusion, the bound for the second Hamming weight based on the second Feng-Rao number is better than those given by Kirfel-Pellikaan in [KP] and the Griesmer order bound introduced in [DFGL2], for AG codes coming from both generalized Hermitian curves and Suzuki curves. Additionally, recall that this bound equals the one given by the actual second Feng-Rao distance in most cases and specifically for all $a \geq c$.

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